

# Decentralized Systems Engineering

CS-438 – Fall 2024

DEDIS

Pierluca Borsò-Tan and Bryan Ford

EPFL

Credits: P. Tennage, C. Basescu, et al.

# Decentralized Communication

Gossip

(Homework 1)

# Recap: Improved Gossiping

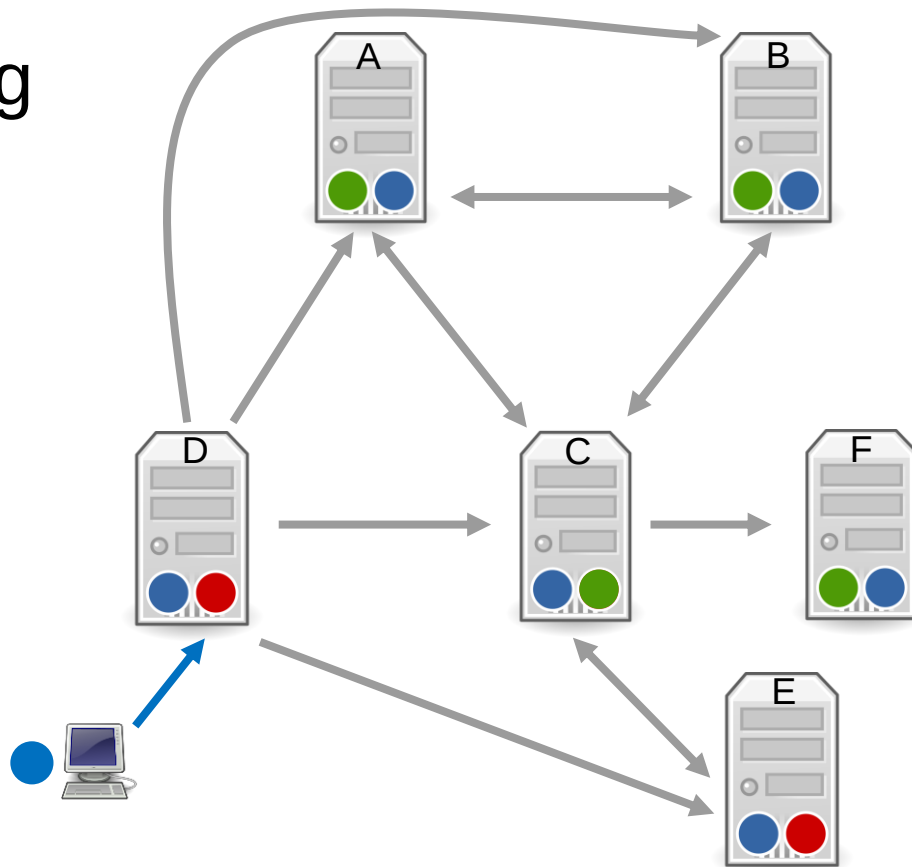
What can we learn from people ?

## Rumor mongering

On receiving message M:

```
pick random neighbor, send M  
neighbor replies: new rumor?
```

```
if new:      repeat  
else:        flip_coin()  
              if head: repeat  
              else:   stop
```



What's good about this ? Which issues do you foresee ?

# Recap: Improved Gossiping (cont'd)

How can make sure messages reach *every* node?

## **Anti-entropy**

Periodically (when timer fires):

```
pick random neighbor  
send "anything new?"  
reduce entropy (ihave/sendme)
```

What's good about this ? What's limiting ?

This (slowly) ensures complete dissemination, at  $O(n)$  per period

# Case study: Twitter

Gossip relies on message-IDs.

How do you pick a good Message-ID in a (centralized) distributed setting ?

- For most applications, they need to be sortable
- Globally unique, even in distributed settings

→ Snowflake IDs

[ 0 (1 bit) | timestamp (41 bits) | machine ID (10 bits) | counter (12 bits) ]

Announced 2010, adopted by Instagram (2012), Discord (2015)

# Case study: Mastodon

Snowflake IDs (64 bits)

[ 0 (1 bit) | timestamp (41 bits) | machine ID (10 bits) | counter (12 bits) ]

What about a *decentralized* setting?

- Would this work?
- What else do we need to do ?

# Case study: Mastodon

```
{
```

```
'created_at': datetime.datetime(2022, 11, 13, 0, 52, 37, tzinfo=tzutc()),
```

```
'id': 109347716491680514,
```



Local Snowflake ID

```
... message data ...
```

```
'uri': 'https://example.com/@user/109347716173491502'
```



Original URI w/ Snowflake ID

```
}
```

# Gossip Quality Measures

Is your gossip any good ?

How can you tell ?

- Residue
- Traffic  $\frac{\text{total traffic}}{\text{number of nodes}}$
- $t_{avg}$  ,  $t_{95}$  average & 95<sup>th</sup> percentile time  
for rumor  $\rightarrow$  node
- $t_{last}$  time for rumor  $\rightarrow$  last node

| Broadcast | Rumor-Mongering          | Anti-Entropy             |
|-----------|--------------------------|--------------------------|
|           | Early: low<br>Late: high | Early: high<br>Late: low |
|           |                          |                          |
|           |                          |                          |
|           |                          |                          |

# What are the parameters & trade-offs ?

Rumor-mongering (fast, randomized):

- Feedback vs. blind
- Randomized vs. counter
- Constant: probability / counter

Anti-entropy (slow, complete):

- Periodicity

# How can we “delete” rumors ?

## Death Certificates

- When is the rumor deleted ?
- How does this affect propagation ?
- Why can rumors resurrect ?
- When can we delete the death certificate ?
- What if a server is offline ?
  - “Dormant” death certificates
  - Phased deletion

# A few applications of Gossip

- Metadata propagation
- Failure detection
- Group membership

## Example

- Apache Cassandra
- CockroachDB
- Consul

# Communicating with (many, unknown) peers

You now understand:

- How do we reach unknown peers ?
- How do we eventually reach *every* peer ?
- How can we communicate reliably ?
  - Are they online ?
  - Is the network « stable » ?
- How do they find out about us / our node ?

# Decentralized Search

Finding Data

(Homework 2)

# Finding data among (many, unknown) peers

- Same machine  
build a local index, and/or Linux `locate`, `find`, `grep` commands
- Local networking
- Global networking, centralized trust  
crawling, processing, indexing, then using the (distributed) index...  
Google 😊
- Decentralized  
??? → today's and next week's lecture

# Finding data among (many, unknown) peers

Many open questions:

- Where can the data be found ?
- Does the data even exist ?
- Who knows about it ?
- How do we retrieve it ?

Most importantly:

- What can we assume about the peers and the network ?

# Distributed search algorithms

Two big families:

- Unstructured search
  - Robust to churn, instantly adaptive
    - Today's lecture
- Structured search
  - Much more efficient, many more problems
    - Next week's lecture

# Building Gnutella

Context (1999 – 2008):

- No Spotify, no Netflix
- No BitTorrent
- People still want entertainment

Specifications:

- Search any file, anywhere
- Metadata can be searched
- Complex queries are allowed ((A or B) and C)  
e.g. artist is “Bob Marley” and title contains “birds”

# Standard, basic algorithm

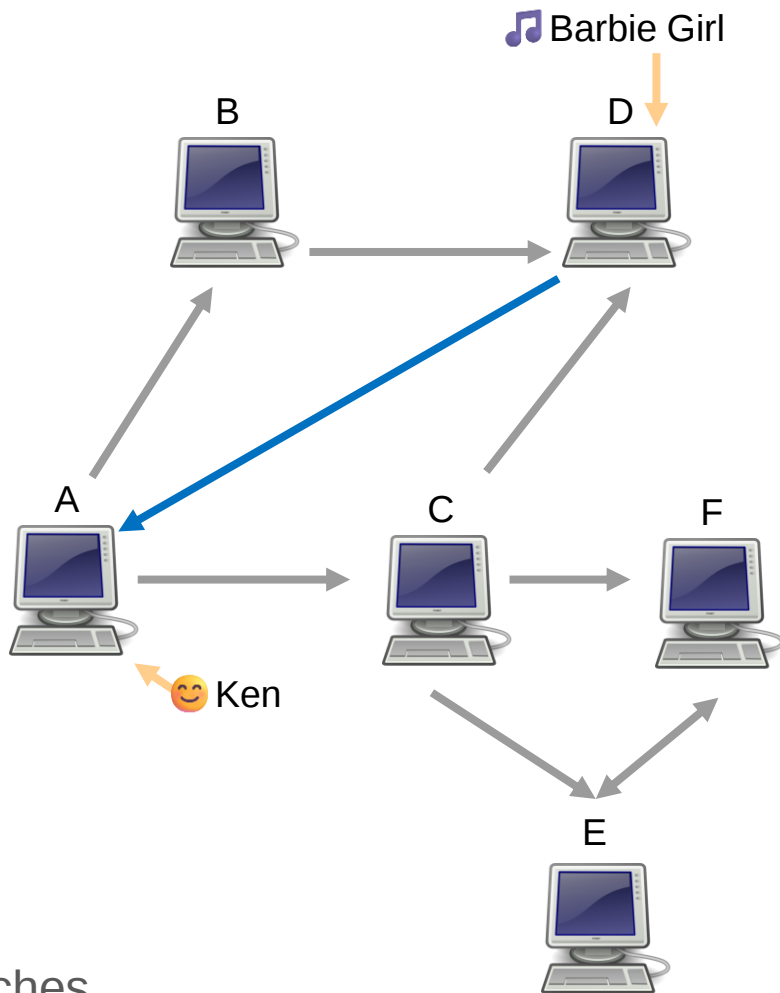
What can we learn from people ?

## Flooding

- Gossip searches (query)
- Direct response (query hit)

Which issues do you foresee ?

- Unpredictable delays
- Connectivity
- Efficiency : all nodes see & process all searches



# Optimizations

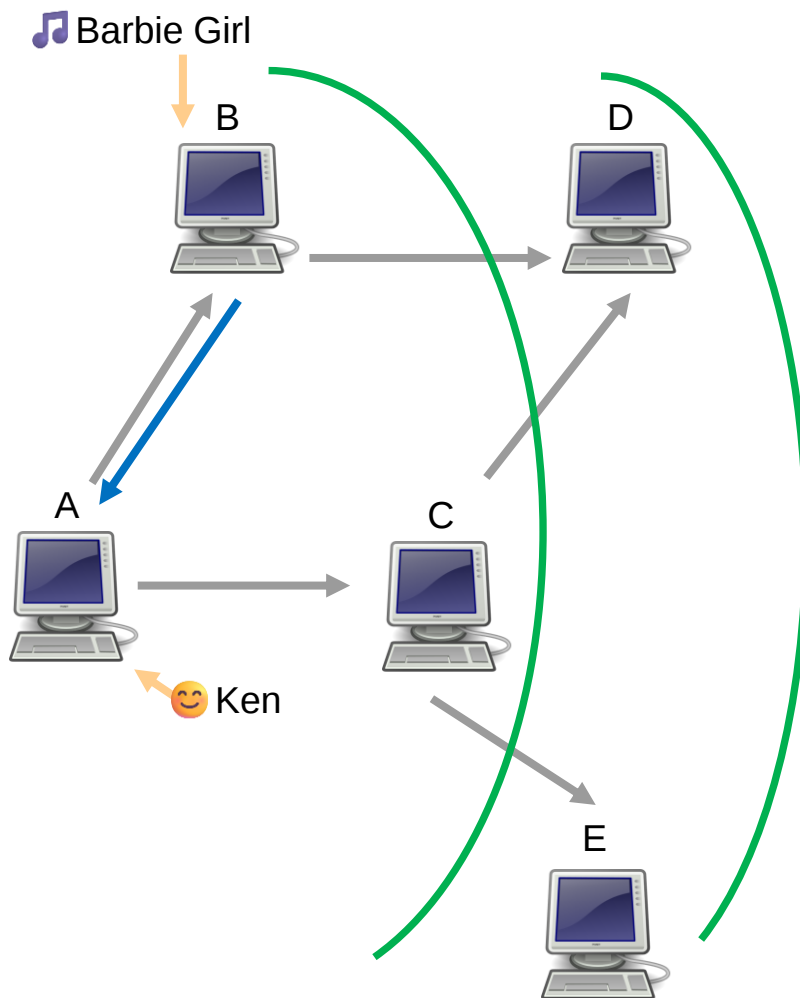
Can we *not* flood everyone ?

## Expanding-ring search

- Limited flooding (TTL)
- Increasing TTL on retry

What are the trade-offs ?

- Higher latency
- Asymptotically worse  $O(n \log n)$
- Pragmatically, it works !



# Optimizations

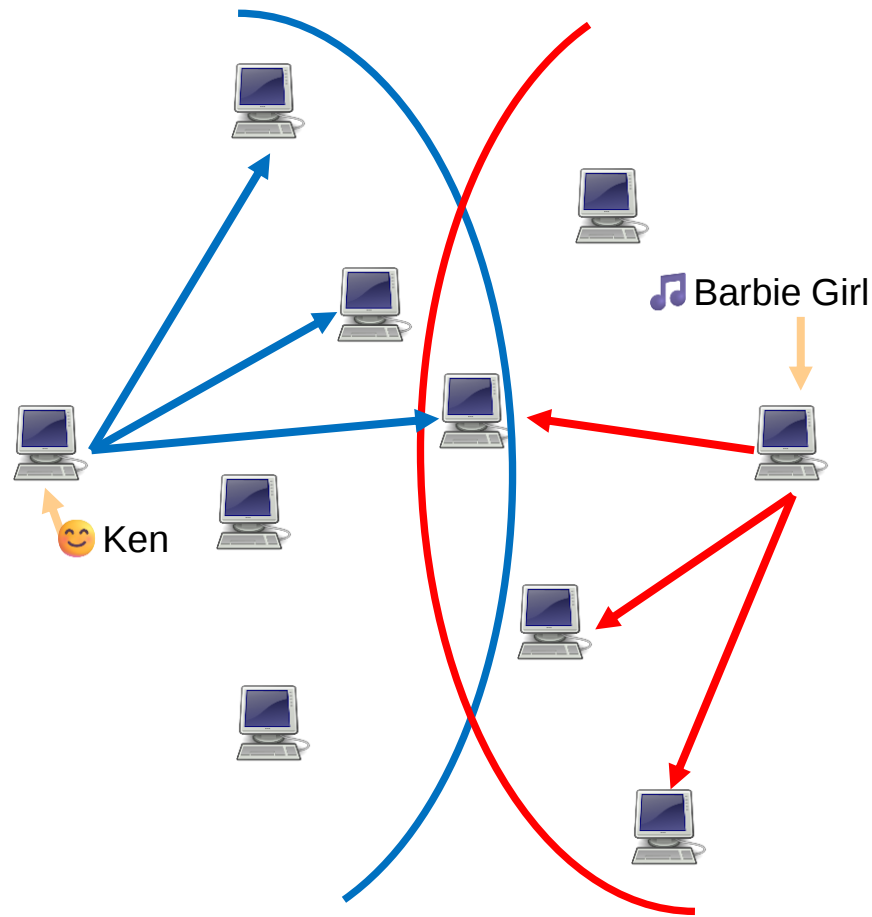
Can we make it more efficient ?

## BubbleStorm

- Birthday paradox
- Data search & storage  
random “meet in the middle”

Key considerations:

- Asymptotically efficient  $O(\sqrt{n})$
- “Mostly unstructured”
- Tunable parameters
- Extremely robust / resilient



# Next steps

Reading on Moodle:

- (Mandatory) BubbleStorm: Resilient, Probabilistic, & Exhaustive P2P Search

→ Use Friday's session to ask questions